

Home Health Care Services: Models and Algorithms

Stefan Nickel, Jörg Steeg, Michael Schröder

Institute for Operations Research
Discrete Optimization and Logistics

Outline

- Motivation and three main problems
- Models and algorithms
- Experiments with real-world data
- Summary and outlook

Motivation

- Aging society
- Decreasing budget per patient
- Shift from hospital care towards home care
- Intensive competition among home health care (HHC) service providers
- People in rural areas live in self-owned houses (which have a small value)

→ Generating improved schedules for HHC services

- Improvement of utilization of resources to reduce costs
- Higher level of comfort for the patients

Motivation (2)

HHC service “We care 4 you”

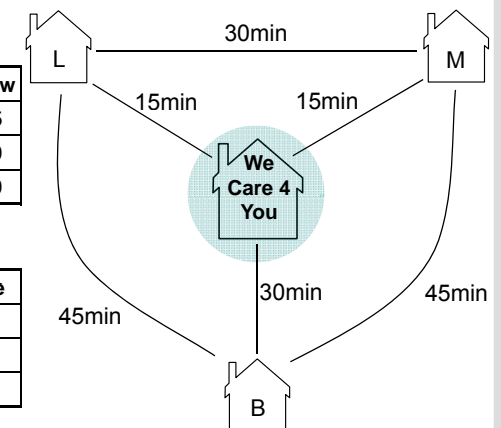
■ Patients

Patient	Service duration	Time window
Bart	30 min	8:00 – 8:45
Liza	30 min	8:30 – 9:30
Maggie	30 min	8:00 – 9:30

■ Distances

	Bart	Liza	Maggie
Bart	0	45	45
Liza	45	0	30
Maggie	45	30	0

■ Nurses: Patty, Selma



Motivation (3)

HHC service “We care 4 you”

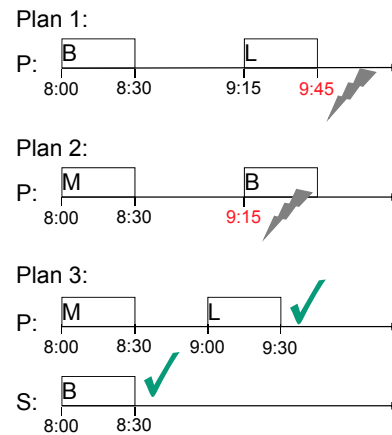
■ Patients

Patient	Service duration	Time window
Bart	30 min	8:00 – 8:45
Liza	30 min	8:30 – 9:30
Maggie	30 min	8:00 – 9:30

■ Distances

	Bart	Liza	Maggie
Bart	0	45	45
Liza	45	0	30
Maggie	45	30	0

■ Nurses: Patty, Selma



Motivation (4)

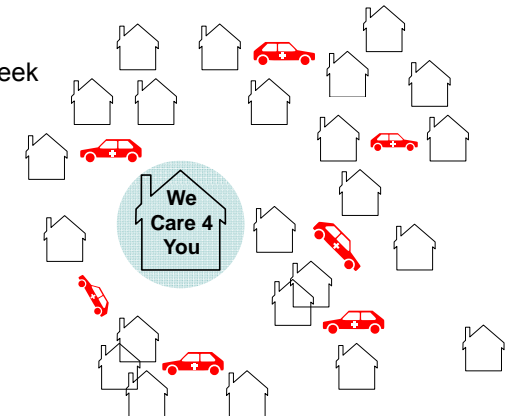
Exemplary HHC service

Patients: 50 = 300 visits per week

- Time windows
- Different shift combinations
- Qualifications

Nurses: 15

- Different availabilities
- Different qualifications



Three main problems

- Home health care problem (HHCP)
 - Feasible assignment of jobs to nurses and start times to jobs
- Master schedule problem (MSP)
 - Mid-term plan as basis for weekly assignment of nurses to jobs
- Operational planning problem (OPP)
 - Incorporate upcoming events (perturbations) into an existing schedule

Outline

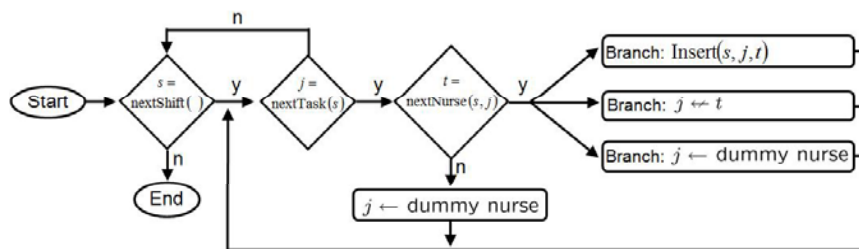
- Motivation and three main problems
- **Models and algorithms**
- Experiments with real-world data
- Summary and outlook

Home Health Care Problem (HHCP)

- Combination of two \mathcal{NP} -hard problems
 - Periodic Vehicle routing problem with time windows
 - Nurse rostering problem
- Task: Feasible assignment of jobs to nurses and start times to jobs
- Given
 - Planning horizon of $S = \{1, \dots, S\}$ shifts with hard time windows $[0, H]$
 - Set of nurses $\mathcal{N} = \{1, \dots, N\}$
 - Set of jobs $\mathcal{J} = \{1, \dots, J\}$
 - Travel distance between jobs $D = (d_{i,k})_{i,k \in \mathcal{J}}$

- Objective function: Minimize the weighted sum of...
 - ... the number of unscheduled tasks $\min \alpha_1 \cdot DNC$
 - ... patient-nurse loyalty penalty $+ \alpha_2 \cdot \sum_{j=1}^J (|X_j| - 1)$
 - ... overtime costs $+ \alpha_3 \cdot \sum_{n=1}^N \max \left\{ 0, \sum_{s \in S} (end_n^s - start_n^s) - wt_n \right\} \cdot c_n$
 - ... travelling distance $+ \alpha_4 \cdot \sum_{n=1}^N \sum_{s=1}^S dt_n^s$
- Exemplary overview of constraints
 - Exactly one of the possible shift combinations is chosen, and a task representing the job is scheduled in each of those shifts.
 - Each job is performed during its time window.
 - A nurse can only be assigned to a job if she possesses the required qualification.
 - Each route starts and ends at the depot within the shift's time window

- Solution approach: Two-stage algorithm
- I. Constraint programming heuristic



→ Quickly generate a feasible solution

II. Adaptive large neighborhood search (ALNS)

1. function ALNS(solution $s, q \in \mathbb{N}, \sigma \in \mathbb{N}$)
2. Solution $s_{best} = s$
3. $m = 0$
4. repeat
5. $m = m + 1$
6. $s' = s$
7. Choose removal operation r with probability π_r
8. Choose insertion operation i with probability p_i
9. Remove q requests from s' with operation r
10. Reinsert removed requests into s' with operation i
11. if $f(s') < f(s_{best})$ then
12. $s_{best} = s'$
13. if $\text{accept}(s', s)$ then
14. $s = s'$
15. if $(m \bmod \sigma) == 0$ then
16. Update probabilities for operations
17. until (stopping-criterion is met)
18. return s_{best}

→ CP layer maintains feasibility while improving the solution

Models and algorithms – HHCP



- Removal heuristics
 - Random removal
 - Shift combination removal
 - Worst removal
- Insertion heuristics
 - In order insertion
 - Greedy insertion
- Score for each insertion and removal heuristic that stores how successful the heuristic was in the past

→ Preference of previously successful operations

Models and algorithms – MSP



Master schedule problem (MSP)

- Task: Mid-term plan as basis for weekly assignment of nurses to jobs
- Given
 - Planning horizon of $S = \{1, \dots, S\}$ shifts with hard time windows $[0, H]$
 - Set of jobs $\mathcal{J} = \{1, \dots, J\}$
 - Travel distance between jobs $D = (d_{i,k})_{i,k \in \mathcal{J}}$
 - Maximal tour length L
- Objective function: Minimize the number of tours required to visit all patients

Models and algorithms – MSP

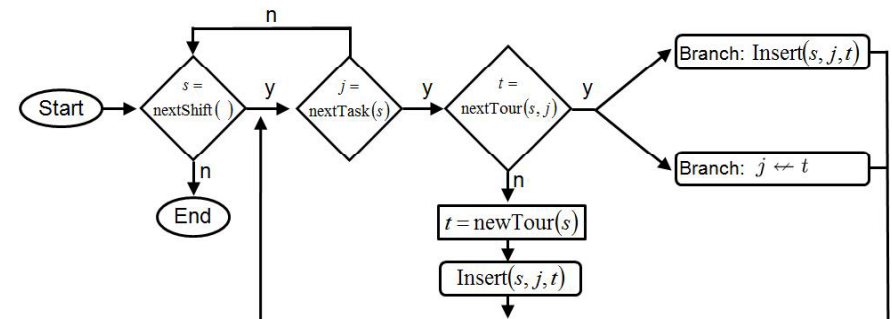


- Exemplary overview of constraints
 - Each job must be performed as often as indicated by its frequency
 - Each task is assigned to exactly one tour in each shift
 - Each tour has a maximal length
 - A tour only contains tasks of one \mathcal{L}_j qualification dimension
 - Each tour starts and ends at the depot within the shift's time window
- Jobs only have one possible shift combination
- Solution approach:
 - Constraint programming heuristic
 - Generate lower and upper bounds on the minimal number of tours

Models and algorithms – MSP



- Constraint programming heuristic



→ Construct tours such that next task is scheduled at closest tour that is compatible with required qualification

Operational planning problem (OPP)

- Task: Incorporate upcoming events (perturbations) into existing schedule
 - Non-availability of personnel
 - Changes in jobs
 - **Incorporation of a new patient**
- Indicators to measure perturbation: Inconvenience for...

- ... nurses
- ... patients

$$\varphi_j^s = \begin{cases} 0 & , \text{ if } x_j^s = 0, \\ |st_j^s - \widehat{st}_j^s| & , \text{ if } x_j^s \neq 0, \widehat{x}_j^s \neq 0, \\ 0 & , \text{ if } x_j^s \neq 0, \widehat{x}_j^s = 0 \text{ and } st_j^s \in [hbs_j, hbe_j - pt_j] \\ \min \left\{ \begin{array}{l} |st_j^s - hbs_j|, \\ |st_j^s + pt_j - hbe_j| \end{array} \right\} & , \text{ if } x_j^s \neq 0, \widehat{x}_j^s = 0 \text{ and } st_j^s \notin [hbs_j, hbe_j - pt_j]. \end{cases}$$

→ Deviation of start times of initial to new solution

- Objective function: Minimize...
 - ... perturbation penalty φ_j^s
 - ... the original objectives of the HHCP $\widehat{\zeta}$

$$\min \beta_1 \cdot \widehat{\zeta} + \beta_2 \cdot \sum_{s \in \mathcal{S}} \sum_{j \in \mathcal{J}} \varphi_j^s$$

- Solution approach: Two-stage algorithm
 - I. Constraint programming (CP) insertion heuristic**
 - Time window of the new task ends before the nurse-shift starts
 - Nurse-shift ends before the time window of the new task starts
 - Time window of the new task (partly) overlaps the nurse-shift

→ Quickly generate a feasible solution

II. Tabu search algorithm

- 1: **Input:** OPP Θ , Initial solution σ , Move-List $\mathcal{M} = \{1, \dots, M\}$
- 2: $k = 0, T = \emptyset, \sigma^* = \sigma$
- 3: **while** termination criterion not satisfied **do**
- 4: $k = k + 1$
- 5: Select $m \in \mathcal{M}$
- 6: **if** $N_m(\sigma) \setminus T \neq \emptyset$ **then**
- 7: $\sigma_k = \operatorname{argmin}_{s \in N_m(\sigma)} \{\zeta(s)\}$
- 8: **if** $\zeta(\sigma_k) < \zeta(\sigma^*)$ **then**
- 9: $\sigma^* = \sigma_k$
- 10: **end if**
- 11: $\sigma = \sigma_k$
- 12: Update T mit σ_k
- 13: **end if**
- 14: **end while**

→ Improve the solution until a termination criterion (time limit or move limit) is met

→ CP layer maintains feasibility while improving the solution

- Motivation and three main problems
- Models and algorithms
- **Experiments on real-world data**
- Summary and outlook

Experiments on real-world data

- Two test instances
 - I. Health care service provider in a rural area in Germany
 - Long traveling distances up to 45 minutes
 - Many short but also several long service times (> 3 hours)
 - II. Health care service provider in an urban area in the Netherlands
 - Short traveling distances
 - Only short service times (≤ 45 minutes)

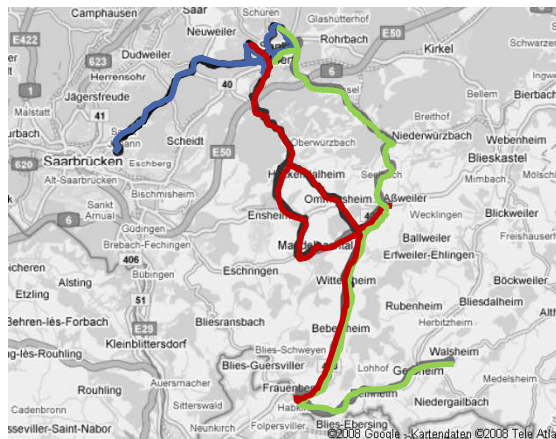
Instance	Nurses	Patients	Tasks
D	11	53	287
NL	12	95	361

Experiments on real-world data

- Solution created by commercial software
 - 55 unscheduled tasks (assigned to dummy nurse)
 - 0 min of overtime (instance taken as default for working times)
 - Solution created by the ALNS for the HCCP
 - 19 unscheduled tasks (assigned to dummy nurse)
 - 0 min of overtime
- Adaption of the nurses' time windows leads to an improved schedule
- Termination criterion: 1000 moves
 - Initial solution: 11 sec

Experiments on real-world data

- Example of three routes created by the ALNS



Experiments on real-world data

- Solution created by the ALNS for the HHCP
 - 45 nurse shifts
 - Solution created by the CP heuristic for the MSP
 - 53 tours
- Negotiation of time windows and nurse shifts can reduce the number of tours
- Computation time: 19.5 sec

Outline



- Motivation and three main problems
- Models and algorithms
- Experiments on real-world data
- **Summary and outlook**

Summary



- Set of algorithms to support the planning process of HHC services
 - Complete re-creation of a schedule
 - Creation of a master schedule and incorporation of unforeseen events
- Introduction of shift combinations for jobs
- Definition of patient-nurse loyalty: Minimize the number of different nurses per patient throughout the schedule
- Development of hybrid CP-meta-heuristic algorithms
- Development of new operators and a new application for ALNS
- Proved applicability to real-world data sets

Outlook: Further research topics



- Home health care problem
 - Introduction of split services
 - Implementation of further propagation techniques (whenDomain)
- Master schedule problem
 - Minimization of qualification level for a tour
 - Introduction of a second target length for a route
- Operational planning problem
 - Development of a decision support system
 - Incorporation of further events into the OPP



Home Health Care Services: Models and Algorithms

Stefan Nickel, Jörg Steeg, Michael Schröder

Institute for Operations Research
Discrete Optimization and Logistics

Thank you for your attention!

http://www.itwm.fhg.de/zentral/download/berichte/bericht_173.pdf

Home health care problem (HHCP)

- E. Cheng and J. L. Rich. A Home Health Care Routing and Scheduling Problem. Technical Report TR98-04, Department of CAAM, Rice University, USA, 1998.
- S. Bertels and T. Fahle. A hybrid setup for a hybrid scenario: combining heuristics for the home health care problem. *Computers & Operations Research*, 33:2866-2890, 2006.
- P. Ekebom, P. Flisberg, and M. Rönnqvist. LAPS-CARE – an operational system for staff planning of home care. *European Journal of Operational Research*, 171:962-976, 2006.

Master schedule problem (MSP)

- J. Desrosiers, M. Sauvé, and F. Soumis. Lagrangian relaxation methods for solving the minimum fleet size multiple travelling salesman problem with time windows. *Management Science*, 34(8):1005-1022, August 1988.
- M. Surico, U. Kaymak, D. Naso, and R. Dekker. Hybrid Meta-Heuristics for Robust Scheduling. Technical report, Erasmus University Rotterdam, Netherlands, 2006.
- N. Azi, M. Gendreau, and J.-Y. Potvin. An Exact Algorithm for a Vehicle Routing Problem with Time Windows and Multiple Use Of Vehicles. Presented at the „Triennial Symposium on Transportation Analysis (TRISTAN VI)“, Phuket, Thailand, June 2007.

Operational planning problem (OPP)

- R. Barták, T. Müller, and H. Rudová. Minimal Perturbation Problem – A Formal View. *Neural Network World*, 13(5):501-511, 2003.
- H. E. Sakkout, T. Richards, and M. G. Wallace. Minimal Perturbation in Dynamic Scheduling. In *13th European Conference on Artificial Intelligence, Brighton, UK*, pages 504-508. John Wiley and Sons, 1998.
- J.-F. Cordeau, G. Laporte, and A. Mercier. A unified tabu search heuristic for vehicle routing problems with time windows. *Journal of Operational Research Society*, 52:928-936, 2001.

Example

- Constraint: Exactly one of the possible shift combinations is chosen

\forall Jobs j :

$$\text{element}(\rho_j, r_j, R_{j,1}, \dots, R_{j,K_j}) = \{(e, f, d_1, \dots, d_n) : e \in D(\rho_j), f \in D(r_j), \forall i : d_i \in D(R_{j,i}), f = d_e\}$$

- Propagation:

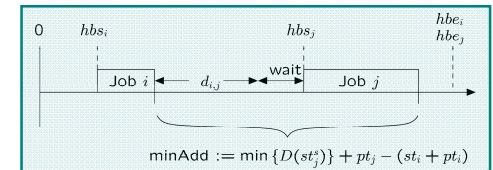
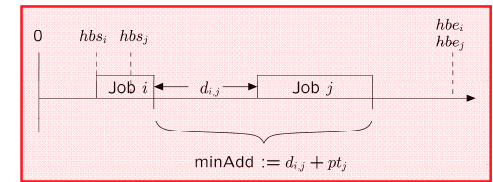
$$\forall s \in r_j : z_j^s = 1 \quad \forall t \notin r_j : z_j^t = 0$$

$$D(x_j^s) = D(x_j^s) \setminus \{-1\} \quad x_j^t = -1$$

$$st_j^t = \min \{D(st_j^t)\}$$

- Algorithm getNextTour(Shift s , Job j)

- $D = \infty, T = -1$
- for all $t \in \mathcal{T}$
- if $y_t \neq s$ then next t
- if $tq_t \parallel req_j$ then next t
- if $t \notin D(x_s^t)$ then next t
- $i =$ currently last job on tour t
- $\text{minAdd} := \max\{d_{i,j} + pt_j, \min\{D(st_j^s)\} + pt_j - (st_i^s + pt_i)\}$
- if $dur_t + \text{minAdd} \leq L$ then
- if $d_{i,j} < D$ then
- $D = d_{i,j}, T = t$
- else if $d_{i,j} = D$ then
- if $dur_t > dur_T$ then
- $D = d_{i,j}, T = t$
- return T



Computation

- AMD Athlon X2
- 2.1 GHz
- 2 GB RAM
- Screenshots: ILOG Scheduler Viewer 1.0